

## Free Products of Graphs

- Produce locally finite graphs non-discrete aut. gps.
- Includes known constructions - <sup>Cayley graphs of</sup> RAAG where graph has complete components.
  - trees with edge colourings.

Let  $\Gamma_1, \dots, \Gamma_n$  be simple connected graphs

$$\Gamma_i = (V_i, E_i)$$

Choose  $\tilde{u}_0 = (u_1, \dots, u_n) \in V_1 \times \dots \times V_n$

$\Gamma = (V, E)$  is the free product of  $\Gamma_1, \dots, \Gamma_n$

$$V = \bigcup_{e=0}^{\infty} V^{(e)} \quad \text{where } V^{(e)} \text{ defined recursively.}$$

We will also  $u: V \rightarrow V_1 \times \dots \times V_n$  recursively.  
 $u(\tilde{v}) = (u_1(\tilde{v}), u_2(\tilde{v}), \dots, u_n(\tilde{v}))$

$V^{(e)}$  will be a set of strings of vertices in  $V_1 V_2 \dots V_n$ .

$V^{(0)} = \{\emptyset\} \quad u(\emptyset) = \tilde{u}_0$       Suppose that  $V^{(e)}$  has been defined.

$$V^{(e+1)} = \left\{ v_1 \dots v_e v_{e+1} : v_1 \dots v_e \in V^{(e)}; v_e \in V_i, v_{e+1} \in V_j \text{ with } i \neq j; \text{ and } v_{e+1} \neq u_j(v_1 \dots v_e) \right\}$$

$$\text{Define } u_i(v_1 \dots v_e v_{e+1}) = \begin{cases} u_i(v_1 \dots v_e) & \text{if } i \neq j \\ v_{e+1} & \text{if } i = j \end{cases}$$

$$E = \left\{ \{\phi, v_i\} : \{u_j(\phi), v_i\} \in E_j \text{ for some } j \right\}^{v_i \in V_i, v_{i+1} \in V_j}$$

$$\cup \left\{ \{\tilde{v}, \tilde{v}v_{e+1}\} : \{u_j(\tilde{v}), v_{e+1}\} \in E_j, i \neq j \right\}$$

$$\tilde{v} = v_1 \dots v_e$$

$$\cup \left\{ \{v_1 \dots v_{e-1}v_e, v_1 \dots v_{e-1}v'_e\} : \{v_e, v'_e\} \in E_i \right\} \quad (*)$$

Note: In (\*), if  $v'_e = u_i(v_1 \dots v_{e-1})$ , then  $v_1 \dots v_{e-1}v'_e \notin V^{(e)}$ .

We use  $v_1 \dots v_{e-1}$  instead.

The number of neighbours of  $\tilde{v} = v_1 \dots v_e$  is

$$N_{\Gamma}(\tilde{v}) = N_{\Gamma_1}(u_1(\tilde{v})) + N_{\Gamma_2}(u_2(\tilde{v})) + \dots + N_{\Gamma_n}(u_n(\tilde{v}))$$

Prop<sup>n</sup>  $\Gamma$  is connected.

Prop<sup>n</sup>  $u: \Gamma \rightarrow \Gamma_1 \times \dots \times \Gamma_n$  is a surjective graph hm.

For  $v_i \in V^{(i)}$  define

$$\varphi(v_i) = \begin{cases} v'_i v_i, & \text{if } v_i \in V_j, i \neq j \\ v'_1 \dots v'_{i-1} v_i, & \text{if } v_i \in V_i \text{ and } v_i \neq u_i(v'_1 \dots v'_{i-1}) \\ v'_1 \dots v'_{i-1}, & \text{if } v_i = u_i(v'_1 \dots v'_{i-1}) \end{cases}$$

$$\varphi(v_i) \in V' \quad \text{and} \quad u(\varphi(v_i)) = u(v_i)$$

Repeat.

$\varphi$  is onto      Let  $\tilde{w}' \in V'$ ,       $\tilde{w}' = w'_1 \dots w'_n$

$\varphi$  is one-to-one

Construct  $\varphi^{-1}$  by same argument changing roles of  $\Gamma$  &  $\Gamma'$ .

$\varphi$  is a graph hm

Check. ▮

Prop<sup>n</sup> There is a hm  $\alpha \mapsto \hat{\alpha} : \text{Aut}(\Gamma_i) \rightarrow \text{Aut}(\Gamma)$ .

Prop<sup>n</sup> Let  $\Gamma, \Gamma'$  be defined as above with initial values  $\tilde{u}_0, \tilde{u}'_0$  resp. Then  $\Gamma \cong \Gamma'$ .

Proof: Define  $\varphi: \Gamma \rightarrow \Gamma'$  recursively.

Since  $\Gamma_1, \dots, \Gamma_n$ , there is  $\tilde{v}' \in V(\Gamma')$  with  $u(\tilde{v}') = \tilde{u}_0$ .

Suppose that  $\tilde{v}' = v'_1 \dots v'_e$  where  $v'_e \in V_i$

Put  $\varphi(\phi) = \tilde{v}'$ .

If  $m \geq l$  and  $\tilde{w}' = v_1' \dots v_l' w_{l+1}' \dots w_m'$ ,

then  $\tilde{w}' = \varphi(w_{l+1}' \dots w_m')$ , where  $w_{l+1}' \dots w_m' \in V$ .

Suppose that  $\tilde{w}' = v_1' \dots v_{l-1}' w_l' \dots w_m'$

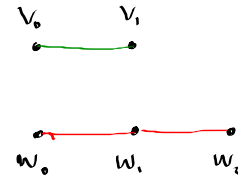
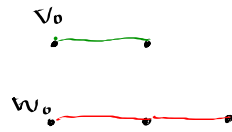
with  $w_{l+1}' \dots w_m' = \emptyset$ ,  $w_l' \neq v_l'$ .

In the first case  $\varphi(u_i(v_1' \dots v_{l-1}')) = w'$

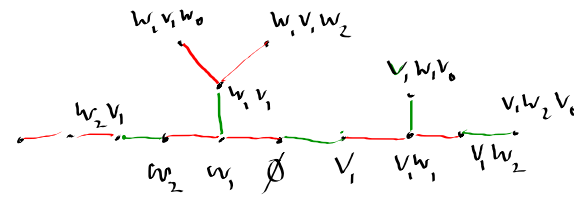
In the second case, if  $w_l' \notin V_0$ , then

$$w' = \varphi(u_i(v_1' \dots v_{l-1}') w_l' \dots w_m')$$



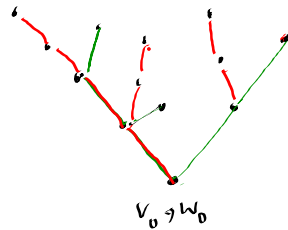


$$\tilde{A}_0 = (V_0, W_0)$$

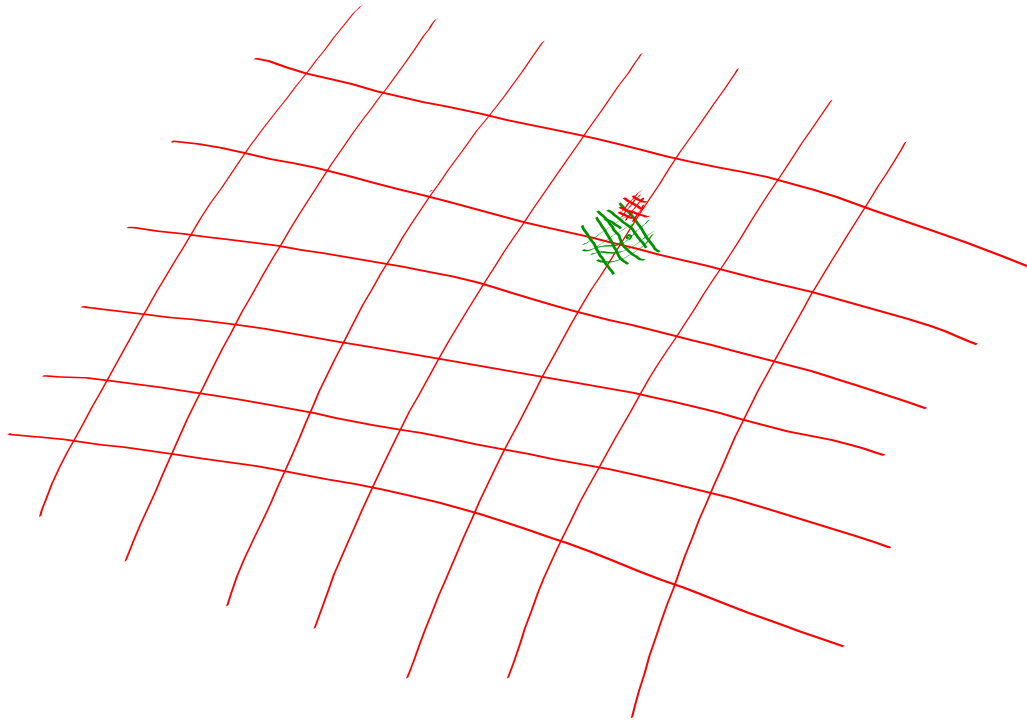


$$u(W_1) = (V_0, W_1) \quad u(V_1) = (V_1, W_0)$$

$$u(W_2) = (V_0, W_2)$$



$\Gamma_1$    $\Gamma_2$  



$\downarrow$   $\downarrow$

